## Modules in model theory

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## 1 pp-definable groups

We study the left modules over a given associative ring R with identity (we do not require commutativity). From our point of view an R-module will be a structure over the language  $L_R = \{0, +, -, r\}_{r \in R}$ , so that a module is effectively an abelian group endowed with a family of endomorphism for every element of R.

**Notation** In the following text x, y, z will denote single variables,  $\boldsymbol{x}$  will denote a tuple of variables  $x_1, \ldots, x_n$  and in this case we define  $|\boldsymbol{x}| := n$  to be the length of the tuple.  $r, r_1, \ldots$  will denote elements of R.

**Definition** We call *equation* an atomic formula:

 $r_1x_1 + r_2x_2 + \dots + r_nx_n = 0$ 

and positive primitive formula (ppf) a formula of type:

$$\exists \boldsymbol{z} \ \gamma_1(\boldsymbol{x}, \boldsymbol{z}) \wedge \cdots \wedge \gamma_n(\boldsymbol{x}, \boldsymbol{z})$$

where the  $\gamma_i$  are equations.

The concept of pp-formula is most important, so we would like to give an alternative interpretation. Suppose we have the pp-formula:

$$\varphi(\boldsymbol{x}) \equiv \exists \boldsymbol{z} \ \gamma_1 \wedge \dots \wedge \gamma_n$$

given x we can look at it as a proposition about the existence of a solution z to a system of equations, or, alternatively, we ask if for a given vector x is there a solution z to the equation:

$$A\boldsymbol{z} = B\boldsymbol{x}$$

where  $A \in B$  are matrices with coefficients in R.

Before moving on with the theory, let's look at some of examples and some properties of pp-formula:

**Example** Suppose R = k is a field and  $M = {}_{k}k$ . We want to study the set defined by the formula  $\varphi(\boldsymbol{x})$  with  $\boldsymbol{x} = (x_1, \ldots, x_n)$ . As we said this is the set of vectors  $\boldsymbol{x}$  for which the system  $A\boldsymbol{z} = B\boldsymbol{x}$  has a solution  $\boldsymbol{z}$ . By a change of basis (Gauss) we can rewrite it as:

$$\begin{pmatrix} I & 0 \\ 0 & 0 \end{pmatrix} \begin{pmatrix} \boldsymbol{z}_1 \\ \boldsymbol{z}_2 \end{pmatrix} = \begin{pmatrix} B'_{11} & B'_{12} \\ B'_{21} & B'_{22} \end{pmatrix} \begin{pmatrix} \boldsymbol{x}_1 \\ \boldsymbol{x}_2 \end{pmatrix}$$

We now see that the set defined by the equation is:

$$\varphi(k^n) = \ker \begin{pmatrix} B'_{21} & B'_{22} \end{pmatrix}$$

Note that this is not the same as considering  $M = k^n$ , in fact in this case the only pp-definable set turns out to be M and 0.

**Example** Let  $\varphi(x)$  be as in the previous example, but suppose now that R is a PID. This time we can't use Gauss reduction, but we can still use Smith normal form to rewrite the equation as:

$$D\boldsymbol{z} = B'\boldsymbol{x},$$

where D is a diagonal matrix. This means that the formula  $\varphi(\boldsymbol{x})$  is equivalent to  $\varphi'(\boldsymbol{x}) \equiv \exists \boldsymbol{z} \ \gamma'_1 \wedge \cdots \wedge \gamma'_n$  where each  $\gamma'_i(\boldsymbol{x}, \boldsymbol{z})$  is of type:

$$d_i z_i = b'_{i1} x_1 + \dots + b'_{in} x_n$$

We can also take Smith normal form of B and rewrite the equation as

$$A'\boldsymbol{z} = \tilde{D}\boldsymbol{x},$$

where  $\tilde{D}$  is a diagonal matrix. In this case each  $\gamma'_i(x, z)$  is of type:

$$a'_{i1}z_1 + \dots + a'_{in}z_n = d_i x_i$$

In general we can't give a more explicit description of a pp-definable set. Still we can prove some important properties.

**Proposition 1.1.** Let  $\varphi(x_1, \ldots, x_n)$  be a pp-formula. The set  $\varphi(M^n)$  is a subgroup of  $M^n$ . If moreover R is commutative then it is a submodule.

*Proof.* Let  $A\mathbf{z} = B\mathbf{x}$  be the equation associated with  $\varphi$ . The zero is in  $\varphi(M^n)$ , because the equation  $A\mathbf{z} = B\mathbf{0}$  always has the trivial solution  $\mathbf{z} = \mathbf{0}$ . Let now  $\mathbf{x}_1$  and  $\mathbf{x}_2$  be in  $\varphi(M^n)$ . This means that we can find  $\mathbf{z}_1$  and  $\mathbf{z}_2$  such that  $A\mathbf{z}_1 = B\mathbf{x}_1$  and  $A\mathbf{z}_2 = B\mathbf{x}_2$ . The vector  $\mathbf{x}_1 - \mathbf{x}_2$  is then in  $\varphi(M^n)$ , because the equation  $A\mathbf{z} = B(\mathbf{x}_1 - \mathbf{x}_2)$  has a solution (take  $\mathbf{z} = \mathbf{z}_1 - \mathbf{z}_2$ ). The last point follow immediately from  $Ar\mathbf{z} = rA\mathbf{z} = r\mathbf{x}$ , for any  $r \in R$ .

We can easily see that, for a pp-formula  $\varphi(\boldsymbol{x}, \boldsymbol{y})$ , the formula  $\varphi(\boldsymbol{x}, \boldsymbol{0})$  still defines a group. The following proposition gives a characterization of the set defined by  $\varphi(\boldsymbol{x}, \boldsymbol{a})$ .

**Proposition 1.2.** Let  $\varphi(\mathbf{x}, \mathbf{y})$  be a pp-formula and  $\mathbf{a} = (a_1, \ldots, a_m)$  be a sequence of elements in M. Then the set  $\varphi(M^n, \mathbf{a})$  is empty or a coset of  $\varphi(M^n, \mathbf{0})$ .

*Proof.* If  $\varphi(M^n, \boldsymbol{a})$  is not empty, fix  $\boldsymbol{x}_0$  in  $\varphi(M^n, \boldsymbol{a})$ . If  $\boldsymbol{x}_1$  is in  $\varphi(M^n, \boldsymbol{0})$  then  $\boldsymbol{x}_0 + \boldsymbol{x}_1$  is in  $\varphi(M^n, \boldsymbol{a})$  because the associated system:

$$A\boldsymbol{z} = B\begin{pmatrix}\boldsymbol{x}_0 + \boldsymbol{x}_1\\\boldsymbol{a}\end{pmatrix}$$

is easily seen to have a solution. On the other hand if  $\boldsymbol{x}_0$  and  $\boldsymbol{x}_1$  are in  $\varphi(M^n, \boldsymbol{a})$  then  $\boldsymbol{x}_1 - \boldsymbol{x}_0$  is in  $\varphi(M^n, \boldsymbol{0})$ .

Last we note that pp-definable subgroups are closed under  $\cap$  and +. In fact we can see any equation  $\gamma(\boldsymbol{x}, \boldsymbol{y})$  in some variables as an equation  $\tilde{\gamma}(\boldsymbol{x}, \boldsymbol{y}, \boldsymbol{z})$  relating more variables, simply attaching zero coefficients to the new variables. Hence we can combine two pp-formulas without mixing up existentials:

$$\begin{array}{lll} (\varphi \cap \psi)(\boldsymbol{x}) &=& \varphi(\boldsymbol{x}) \wedge \psi(\boldsymbol{x}) \\ (\varphi + \psi)(\boldsymbol{x}) &=& \exists \boldsymbol{y}, \boldsymbol{z} \ \varphi(\boldsymbol{y}) \wedge \psi(\boldsymbol{z}) \wedge \boldsymbol{x} = \boldsymbol{y} + \boldsymbol{z}. \end{array}$$

**Example** Let  $M = {}_{R}R$ . Is easy to see that every pp-definable subgroup of M is a right ideal.

Conversely every finitely generated right ideal is pp-definable. In fact let  $g_1, \ldots, g_n$  be the generators of the ideal, then the  $x \in R$  such that

$$\exists \boldsymbol{z} \in R^n \ (g_1, \ldots, g_n) \boldsymbol{z} = x$$

are precisely the elements of the ideal. It follows that every right ideal of a noetherian ring is pp-definable and these are the only pp-definable subgroups. The converse, that is, if every right ideal is pp-definable then R is noetherian, is also true if we assume R to be weakly saturated (the proof is quite simple).

**Example** Let M be an R module. A definable subgroup of M is closed under endomorphism of M. In fact if  $x \in M$  is such that  $\exists z(az = bx)$  then we also have  $\exists z'(az' = b(x\varphi))$ , take  $z' = z\varphi$ .

Let then  $R = \mathbb{Z}$  and  $M = \mathbb{Q}$ . The  $\mathbb{Z}$ -endomorphisms of  $\mathbb{Q}$  act transitively, so by what we just said the only pp-definable subsets of  $\mathbb{Q}$  are 0 and  $\mathbb{Q}$ .

The same is true if we let R = k be a field (or more generally a division algebra) and M a k-vector space.

## 2 Quantifier elimination

We want to prove the following weak form of quantifier elimination.

**Theorem 2.1.** For every module M, every  $L_R$ -formula is equivalent to a boolean combination of positive primitive formulas. That is, given a formula  $\psi(\mathbf{x})$  we can find  $\varphi(\mathbf{x})$  a boolean combination of pp-formulas so that:

$$M \models \psi(\boldsymbol{x}) \leftrightarrow \varphi(\boldsymbol{x})$$

for every  $\boldsymbol{x}$  in  $M^n$ .

Let's first introduce some convenient terminology. Fix a group G. We say that a subset X of G is G-big if a finite number of translations of X cover G, else we say that X is G-small. Note that a subgroup H of G is G-big if and only if G/H is finite. We leave to the reader to verify that a finite union of small sets is small.

**Lemma 2.2** (B.H. Neumann). Let  $H_i$  be subgroups of an abelian group G. If  $H_0 + a_0 \subset \bigcup_{i=1}^n H_i + a_i$  and  $H_i \cap H_0$  is small in  $H_0$  for i > k, then  $H_0 + a_0 \subset \bigcup_{i=1}^k H_i + a_i$ 

Proof. Translating everything by  $-a_0$  and taking the intersection with  $H_0$ , the hypothesis reads  $H_0 = \bigcup_{i=1}^n H_i + a_i$  with  $H_i \subset H_0$  and  $H_i$  is  $H_0$ -small for i > k. We must prove that we can throw away the small set. Let  $C = H_0 \setminus \bigcup_{i=1}^k H_i + a_i$ . If C is empty we have finished. If it is not empty then C is necessarily  $H_0$ -big. In fact  $H_1, \ldots, H_k$  are  $H_0$ -big (e.g.  $H_0/H_i$  is finite) and by basic group theory we deduce that  $H_1 \cap \cdots \cap H_k$  too is  $H_0$ -big. Let now c be an element in C, then  $(H_1 \cap \cdots \cap H_k) + c$  is G-big and is contained in C, because  $(\bigcap_{i=1}^k H_j + c) \cap (H_i + a_i) \subset (H_i + c) \cap (H_i + a_i) = \emptyset$  for  $i \leq k$ . But by hypothesis  $C \subset \bigcup_{i=k+1}^n H_i + a_i$  and the latter is a finite union of small set, so it can't contain a big set.

**Lemma 2.3.** Let  $A_i$  be sets. If  $A_0$  is finite, then  $A_0 \subset \bigcup_{i=1}^k A_i$  iff

$$\sum_{\Delta \subset \{1,\dots,k\}} (-1)^{|\Delta|} \left| A_0 \cap \bigcap_{i \in \Delta} A_i \right| = 0.$$

*Proof.* A simple application of the inclusion-exclusion principle.

We are now ready to prove the theorem.

Proof of Theorem 2.1. The only thing we have to prove is that if  $\varphi(x, y)$  is equivalent to a boolean combination of pp-formulas, so is  $\psi(y) \equiv \forall x \varphi(x, y)$ . Note that pp-formulas are closed under conjunction, so we can write:

$$\varphi \equiv \neg \varphi_0 \lor \varphi_1 \lor \cdots \lor \varphi_k \equiv \varphi_0 \to \varphi_1 \lor \cdots \lor \varphi_k$$

where  $\varphi_i$  are pp-formulas. Set-wise this means that  $M \models \psi(\boldsymbol{y})$  iff  $\varphi_0(M, \boldsymbol{y}) \subset \varphi_1(M, \boldsymbol{y}) \cup \cdots \cup \varphi_k(M, \boldsymbol{y})$ . Setting  $H_i = \varphi_i(M, 0)$ , by Proposition 1.2 we can rewrite this as  $H_0 + a_0 \subset \bigcup_{i=1}^n H_i + a_i$  for some  $a_i$  in  $M^n$ . By Lemma 2.2 we can assume  $H_0 + a_0 \subset \bigcup_{i=1}^k H_i + a_i$  and  $H_0/H_i \cap H_0$  finite. We now have to find a boolean combination of pp-formulas that express this inclusion, but being  $H_0$  infinite this isn't a simple task (if it were finite we could simply impose the inclusion element by element). However if we take the quotient by  $H_0 \cap \cdots \cap H_k$ (a  $H_0$ -big set) we are left with the inclusion of a finite set:

$$H_0/(H_0 \cap \dots \cap H_k) + a_0 \subseteq \bigcup_{i=1}^k H_i/(H_0 \cap \dots \cap H_k) + a_i$$
(1)

We can now apply Lemma 2.3 to (1). Let  $N_{\Delta}$  be

$$N_{\Delta} = \left| \left( H_0 \cap \bigcap_{i \in \Delta} H_i \right) / (H_0 \cap \dots \cap H_k) \right|.$$

The set  $((H_0 + a_0) \cap \bigcap_{i \in \Delta} (H_i + a_i))/(H_0 \cap \cdots \cap H_k)$  is empty or it has  $N_{\Delta}$  elements (Proposition 1.2), so Lemma 2.3 reads:

$$M \models \forall x \ \varphi \ \Leftrightarrow \ \sum_{\Delta \in \mathcal{N}} (-1)^{|\Delta|} N_{\Delta} = 0, \tag{2}$$

where

$$\mathcal{N} = \left\{ \Delta \subset \{1, \dots, k\} \mid \exists x \left( \varphi_0(x, \boldsymbol{y}) \land \bigwedge_{i \in \Delta} \varphi_i(x, \boldsymbol{y}) \right) \right\}$$

We have to prove that the sum in (2) can be written as a boolean combination of pp-formulas. To do this, list all the (finite)  $\mathcal{N}$  for which the sum is zero and write a formula that says that we are in one of those cases. It is easily seen that this can be done with boolean combination of pp-formulas.

**Corollary 2.4.** Two *R*-modules  $M_1$  and  $M_2$  are elementary equivalent iff for every  $ppf \varphi \subseteq \psi$  we have

$$\varphi/\psi(M_1) = \varphi/\psi(M_2),$$

where by  $\varphi/\psi(M)$  we mean  $[\varphi(M):\psi(M)]$  if it is finite, or else  $\infty$ .

## References

 Mike Prest. Model Theory and Modules. London Mathematical Society Lecture Note Series. Cambridge University Press, 1988.