## Structures preserved by the $Q R$-algorithm

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- The shifted $Q R$-algorithm
(2) Polynomial structures
- Definition
- Examples
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- Nonsingular case


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## The shifted $Q R$-algorithm

- Given a matrix $A \in \mathbb{C}^{n \times n}$.

We want to compute the eigenvalues, eigenvectors of $A$. QR-step: given $A^{(\nu)}$, we compute

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- Initialization of the shifted QR-algorithm: $A^{(0)}=A$.

QR-step: given $A^{(\nu)}$, we compute

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\begin{gather*}
A^{(\nu)}-\lambda I=Q R  \tag{1}\\
A^{(\nu+1)}=R Q+\lambda I, \tag{2}
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- (1) and (2) imply the similarity relations

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\begin{aligned}
& A^{(\nu+1)}=Q^{H} A^{(\nu)} Q \\
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- Preservation of structure under the shifted QR-algorithm:
(3) $\Rightarrow$ polynomial structures
(4) $\Rightarrow$ rank structures.


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## Definition

- A polynomial structure on $\mathbb{C}^{n \times n}$ is defined as a collection $\mathcal{P}=\left\{p_{k}\right\}_{k}$, where each $p_{k}$ is a polynomial in 7 variables.
- A matrix $A$ is said to satisfy the structure $\mathcal{P}=\left\{p_{k}\right\}_{k}$ if for every $k$,

for certain
- $\mathrm{Herm}_{k}$ Hermitian
- Uni ${ }_{k}$ unitary,
- $(\mathrm{Rk} r)_{k}$ of rank at most $r$ - $\mathcal{M}$ : set of matrices satisfying $\mathcal{P}$


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p_{k}\left(A, A^{H}, A^{-1}, A^{-H}, \operatorname{Herm}_{k}, \operatorname{Uni}_{k},(\operatorname{Rk} r)_{k}\right)=0
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PROOF.

- Any unitary matrix $Q$ can be 'pulled through' such a polynomial relation:

$$
\begin{aligned}
Q^{H} p(A, & \left.A^{H}, A^{-1}, A^{-H}, \operatorname{Herm}, \mathrm{Uni}, \operatorname{Rk} r\right) Q \\
& =p\left(A_{Q}, A_{Q}^{H}, A_{Q}^{-1}, A_{Q}^{-H}, \operatorname{Herm}_{Q}, \operatorname{Uni}_{Q},(\operatorname{Rk} r)_{Q}\right)
\end{aligned}
$$

where

$$
\begin{array}{r}
A_{Q}:=Q^{H} A Q, \quad \operatorname{Herm}_{Q}:=Q^{H}(\operatorname{Herm}) Q, \quad \operatorname{Uni}_{Q}:=Q^{H}(\mathrm{Uni}) Q, \\
(\operatorname{Rkr})_{Q}:=Q^{H}(\mathrm{Rk} r) Q .
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- Any unitary matrix $Q$ can be 'pulled through'.
- $\Rightarrow$ Polynomial structures satisfied by $A$, must carry over to $A_{Q}=Q^{H} A Q$. And conversely, by applying the same argument to $Q A_{Q} Q^{H}=Q\left(Q^{H} A Q\right) Q^{H}=A$. In particular, this holds for the $Q$-factor of the shifted QR-algorithm, and hence for the matrices $A^{(\nu)}$ and


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## Definition

- We define a rank structure on $\mathbb{C}^{n \times n}$ as a collection $\mathcal{R}=\left\{\mathcal{B}_{k}\right\}_{k}$ where each $\mathcal{B}_{k}$ is a 'structure block'.

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- $i_{k}$ : row index,
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- $\lambda_{k} \in \mathbb{C}$ : shift element.
- A matrix $A \in \mathbb{C}^{n \times n}$ satisfies the structure $\mathcal{R}$ if for every $k$, $\operatorname{Rank} A_{k}\left(i_{k}: n, 1: j_{k}\right) \leq r_{k}, \quad$ where $A_{k}:=A-\lambda_{k} l$.


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## Definition

(Continuation)

- As a special case, $\mathcal{R}$ is called a pure rank structure if all structure blocks $\mathcal{B}_{k}$ have shift element $\lambda_{k}=0$.
- $\mathcal{M}$ : set of matrices which satisfy $\mathcal{R}$ $\mathcal{R}_{\text {pure }}$ : pure rank structure. $\mathcal{M}_{\text {pure }}$ : set of matrices which satisfy $\mathcal{R}_{\text {pure }}$


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## Example

Here is an example of a rank structure $\mathcal{R}=\left\{\mathcal{B}_{1}, \mathcal{B}_{2}\right\}$. The structure block $\mathcal{B}_{1}$ intersects the diagonal and has shift $\lambda_{1}=0.89$, while the structure block $\mathcal{B}_{2}$ is pure:


## Example

Here is an example of a rank structure $\mathcal{R}_{\text {pure }}=\left\{\mathcal{B}_{k}\right\}_{k=1}^{n}$, yielding the class of lower semiseparable matrices:


Allowing shift elements $\lambda_{k}$, we get the class $\mathcal{R}=\left\{\mathcal{B}_{k}\right\}_{k=1}^{n}$ of lower semiseparable plus diagonal matrices.
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## Theorem

(The nonsingular case:) For $A \in \mathcal{M}$ nonsingular we have
(1) rank structure is strictly preserved by applying a $Q R$-step without shift on $A$;
(2) factorizing $A=Q R$, then $Q$ satisfies the pure structure induced by $\mathcal{R}$.
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Theorem
Let $A$ satisfy a structure block $\mathcal{B}_{k}$
By applying a $Q R$-step without shift on $A$, the rank upper bound $r_{k}$ of $\mathcal{B}_{k}$ can increase by at most $\#\left(\mathcal{I}_{\text {dep }, A} \cap \mathcal{I}_{\text {left }, k}\right)$

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$\Rightarrow\left\{\begin{array}{l}\text { Effectively eliminating QR-decompositions } \\ \text { Sparse Givens patterns }\end{array}\right.$

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## Givens transformations

- Given a matrix $A$, we can search a QR-decomposition by solving

$$
\left\{\begin{array}{l}
Q^{H} A=R \\
Q^{H}=\left(G_{n-1, n}^{(n-1)}\right) \ldots\left(G_{2,3}^{(2)} \ldots G_{n-1, n}^{(2)}\right)\left(G_{1,2}^{(1)} \ldots G_{n-1, n}^{(1)}\right)
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$G_{i-1, i}^{(j)}$ : Givens transformation acting on rows $i-1$ and $i$.

- For $n=3$ this specializes to $\left(G_{2,3}^{(2)}\right)\left(G_{1,2}^{(1)} G_{2,3}^{(1)}\right) A=R$ :

$$
\begin{array}{|ccc}
\mathrm{x} & \mathrm{x} & \mathrm{x} \\
\mathrm{x} & \mathrm{x} & \mathrm{x} \\
\mathrm{x} & \mathrm{x} & \mathrm{x}
\end{array} \rightarrow \begin{array}{|ccc}
\mathrm{x} & \mathrm{x} & \mathrm{x} \\
\mathrm{x} & \mathrm{x} & \mathrm{x} \\
0 & \mathrm{x} & \mathrm{x}
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A $Q R$-decomposition $A=Q R$ is called effectively eliminating if each non-trivial $G_{i-1}^{(j)}$; realizes a transition

where $b \neq 0$ lies in the strictly lower triangular part of $A$.

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\left[\begin{array}{cccc}
0 & \ldots & 0 & a \\
0 & \ldots & 0 & b \\
\vdots & & \vdots & \vdots \\
0 & \ldots & 0 & 0
\end{array}\right] \longrightarrow\left[\begin{array}{cccc}
0 & \ldots & 0 & s \\
0 & \ldots & 0 & 0 \\
\vdots & & \vdots & \vdots \\
0 & \ldots & 0 & 0
\end{array}\right]
$$

where $b \neq 0$ lies in the strictly lower triangular part of $A$.

Introduction

## Effectively eliminating QR-decompositions

Example: for the matrix

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A=\left[\begin{array}{ccc}
0 & \times & \times  \tag{5}\\
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we can solve $\left(G_{2,3}^{(2)}\right)\left(G_{1,2}^{(1)} G_{2,3}^{(1)}\right) A=R$ with

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G_{2,3}^{(1)}=G_{1,2}^{(1)}=I_{2}, \quad G_{2,3}^{(2)}=\frac{1}{\sqrt{2}}\left[\begin{array}{rr}
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The effectively eliminating $Q R$-decomposition of $A$ is essentially
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Introduction
Polynomial structures
Rank structures
Singular case

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Introduction

## Outline

(1) Introduction

- The shifted $Q R$-algorithmPolynomial structures
- Definition
- Examples


## Rank structures

- Definition
- Examples
- Nonsingular case

4 Singular case

- Singular case
- Effectively eliminating QR-decompositions
- Sparse Givens patterns


## Sparse Givens patterns

## Definition

Given a pure structure block $\mathcal{B}_{k}=\left(i_{k}, j_{k}, r_{k}\right)$.
We define the staircase shaped set $\mathcal{I}_{\text {Prepare }, k}^{2}$ and the rectangular shaped set $\mathcal{I}_{\text {Skip }, k}^{2}$ as illustrated.


## Sparse Givens patterns

## Reason for introducing $\mathcal{I}_{\text {Prepare }, k}^{2}, \mathcal{I}_{\text {Skip }, k}^{2}$ :

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A $Q R$-decomposition $A=Q R$ is said to satisfy the sparse Givens pattern induced by $\mathcal{R}_{\text {pure }}$ if $G_{i-1, i}^{(j)}=I_{2}$ for all $(i, j) \in \bigcup_{k} \mathcal{I}_{\text {Skip }, k}^{2}$.


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We introduced now sparse Givens pattern induced by any pure structure $\mathcal{R}_{\text {pure }}$.
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When applying a $Q R$-step without shift on $A$, we have the implications
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